

II) Some basics of Differential equations:

So far we have discussed 1st-order system. But we can equivalently consider higher-order equations: (here for simplicity one-parameter)

$$\boxed{\mathcal{L} f(z) = 0} \quad \text{with operator } \boxed{\mathcal{L} = g_n(z) \frac{\partial^n}{\partial z^n} + g_{n-1}(z) \frac{\partial^{n-1}}{\partial z^{n-1}} + \dots + g_0(z)}$$

- * n th-order operator, linear $\rightarrow$ has n -solutions (we take always this order)
- * polynomial coeff. (can multiply with denom. since homog. eq.) (no common factor)
- * $g_n(z)$ also called discriminant locus $\rightarrow$ gives location where operator gets singular (+ possibly ∞)

$\leadsto$ How can we systematically solve these diff. eq's.?

i) An n-th-order operator is equivalent to a system of size n:

Because: $\mathcal{L}f(z) = 0 \Rightarrow \mathcal{D}_z^n f(z) = -\frac{q_0}{q_n} f(z) - \frac{q_1}{q_n} \mathcal{D}_z f(z) - \dots - \frac{q_{n-1}}{q_n} \mathcal{D}_z^{n-1} f(z)$

and we have: $I = (f(z), \mathcal{D}_z f(z), \dots, \mathcal{D}_z^{n-1} f(z))$
length n vector

$\frac{\partial}{\partial z} I = B(z) I$ with $B(z) = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \\ -q_0/q_n & \dots & -q_{n-1}/q_n \end{pmatrix}$ $\leftarrow$ diff. eq. is in last row of connection matrix

and in the other direction we can decouple the system / go to a "derivative basis"

For $B(z)$ generic we have $\mathcal{D}_z J = B(z) J$ with $J = T I$
 $\uparrow$
rational functions inside

Exercise: Find the equivalent 3rd-order operator: (use mathematica)

$B(z) = \begin{pmatrix} 1 & \frac{1}{2} & 0 \\ -\frac{2(5-7z+3z^2)}{3(1-z)^2} & -\frac{4-3z}{3(1-z)} & \frac{2}{3} \\ +\frac{63-1467z-12932z^2+54629z^3-58624z^4+19456z^5}{3(1-z)^3 z^2 (1-16z)^2} & -\frac{9-1035z+18589z^2-32144z^3+19456z^4}{6(1-z)^2 z^2 (1-16z)^2} & -\frac{9-298z+304z^2}{3(1-z)z(1-16z)} \end{pmatrix}$

also find T.

ii) It is also very useful to use the logarithmic derivative / Euler operator

$\Theta := z \frac{\partial}{\partial z}$

Then we have: $\mathcal{L} = \tilde{q}_n(z) \Theta^n + \tilde{q}_{n-1}(z) \Theta^{n-1} + \dots + \tilde{q}_0(z)$

One can prove: $\Theta^n = \sum_{k=1}^n S_2(n, k) z^k \mathcal{D}_z^k$ and $z^n \mathcal{D}_z^n = \prod_{j=0}^{n-1} (\Theta - j)$
(Stirling numbers 2nd kind $S_2(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^i \binom{n}{i} (k-i)^n$)

$\leadsto$ not complicated to go from $\mathcal{D}_z^k$ to Θ^k form of operators

$\leadsto$ see soon why Θ -form is useful

iii) We can classify the singularities of the diff. operator:

z_0 is singular if $q(z_0) = 0$ or possibly $z_0 = \infty$ (use $t=1/z$ to check)

z_0 is a regular singular point if $(z-z_0)^{n-1} q_i(z)/q_n(z)$ are analytic (finite at z_0)

otherwise it is a irregular singular point $\rightarrow$ essential singularity!

Fuchs: at a reg. sing. point there is at least one solution of the form

$$f(z) = (z-z_0)^\alpha \sum_{n=0}^{\infty} a_n z^n$$

α : indicial / local exponent
 $\in \mathbb{C}$ (usually for us $\mathbb{Q}$)

and we can have at z_0 that a solution $f(z)$ is either

analytic, has a pole, or algebraic or logarithmic branch cut

$\rightarrow$ diff. eq. which have only regular sing. points are called FUCHSIAN.

$\rightarrow$ It is expected that FI have always FUCHSIAN diff. eq.s

$\rightarrow$ We call them also PICARD-FUCHS equation

$\wedge$ refers to geometric interpretation of FIs as BAUM-MANU

iv) How do we find all n solutions now? (assume for simplicity $z_0=0$)

First ansatz: $\varpi_0(z) = z^\alpha \sum_{k=0}^{\infty} a_k z^k$
("varpi") $\rightarrow$

Now we find: $\partial z^{k+\alpha} = (k+\alpha) z^{k+\alpha-1}$
because after deriv. we multiply again with z

$\rightarrow$ in ∂ -form it is easy/convenient to work with power series

lowest exponent in z determines α $z=0$

$$\mathcal{L} \varpi_0(z) = \left(\tilde{q}_n(0) \alpha^n + \tilde{q}_{n-1}(0) \alpha^{n-1} + \dots + \tilde{q}_0(0) \right) z^\alpha + z^\alpha \mathcal{O}(z^1)$$

$\stackrel{!}{=} 0$ determines α
indicial equation

*) a_k are determine from higher orders (in ∂ -form it is easy to find a recursion for the coeff.'s a_k or make an ansatz)

*) if indicial equation has n independent solutions we have all solutions found

x) otherwise we add logarithmic solutions of the following type

FROBENIUS
METHOD

$$\boxed{W_1(z) = z^\alpha \left(\log(z) \sum_{k=0}^{\infty} a_k z^k + \sum_{k=0}^{\infty} b_k z^k \right)}$$

and continue with higher log solution until we have found n indep. solutions

$$\boxed{W_k(z) = z^\alpha \left(\frac{1}{k!} \log^k(z) \sum_{i=0}^{\infty} a_i z^i + \frac{1}{(k-1)!} \log^{k-1}(z) \sum_{i=0}^{\infty} b_i z^i + \dots + \sum_{i=0}^{\infty} d_i z^i \right)}$$

- this gives a constructive way to find all solutions
- this gives local solutions valid/convergent at least until the next singularity
- in principle, find basis of solutions around all sing. points and match them in common regions of convergence (analytic continuation) to obtain global solution
- to understand the properties/singularities of a diff. oper. we write down all indicials in a matrix called RIEMANN P-symbol

$$P \left\{ \begin{array}{c|c} z_2 \dots z_k & \\ \hline \alpha_1^{(1)} & \alpha_1^{(n)} \\ \vdots & \vdots \\ \alpha_n^{(1)} & \alpha_n^{(k)} \end{array} \right\}$$

e.g. $P \left\{ \begin{array}{c|c} 0 & 4 & 16 & \infty \\ \hline 0 & 0 & 0 & 1 \\ 0 & 1/2 & 1/2 & 1 \\ 0 & 1 & 1 & 1 \end{array} \right\}$

two square root branch cut

full tower of logarithms

appear $1, \log, \log^2, \dots$

↳ also called a point a maximal unipotent monodromy (MUM) point

→ So in principle we know how to solve n -th d.feq. and therefore also a 1^{st} -order system of size n !

Exercises:

(use mathematica)

Find all local solutions to the diff. operator (up to a truncation order)

$$L^{(3)} = (1-16z)^2 \partial^3 - 48z(1-16z) \partial^2 - 32z(1-24z) \partial - 8z(1-32z)$$

↑
operator of a very
specific special
double box graph

compute a global solution numerically (voluntary).